## CS283：Robotics Spring 2024：Planning

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## Path Planning: Overview of Algorithms

## 1. Optimal Control

- Solves truly optimal solution
- Becomes intractable for even moderately complex as well as nonconvex problems



## 2. Potential Field

- Imposes a mathematical function over the state/configuration space
- Many physical metaphors exist
- Often employed due to its simplicity and similarity to optimal control solutions



## 3. Graph Search

- Identify a set edges between nodes within the free space

- Where to put the nodes?



## Graph Search

- Overview
- Solves a least cost problem between two states on a (directed) graph
- Graph structure is a discrete representation
- Limitations
- State space is discretized $\rightarrow$ completeness is at stake
- Feasibility of paths is often not inherently encoded
- Algorithms
- (Preprocessing steps)
- Breath first
- Depth first
- Dijkstra
- A* and variants
- D* and variants


## Graph Construction: Visibility Graph



- Particularly suitable for polygon-like obstacles
- Shortest path length
- Grow obstacles to avoid collisions


## Graph Construction: Visibility Graph

- Pros
- The found path is optimal because it is the shortest length path
- Implementation simple when obstacles are polygons


## - Cons

- The solution path found by the visibility graph tend to take the robot as close as possible to the obstacles: the common solution is to grow obstacles by more than robot's radius
- Number of edges and nodes increases with the number of polygons
- Thus it can be inefficient in densely populated environments


## Graph Construction: Voronoi Diagram



- Tends to maximize the distance between robot and obstacles


## Topology Graph



## Graph Construction: Voronoi Diagram

- Pros
- Using range sensors like laser or sonar, a robot can navigate along the Voronoi diagram using simple control rules
- Cons
- Because the Voronoi diagram tends to keep the robot as far as possible from obstacles, any short range sensor will be in danger of failing
- Voronoi diagram can change drastically in open areas


## Graph Construction: Exact Cell Decomposition (2/4)



Graph Construction: Approximate Cell Decomposition (3/4)


## Graph Construction: Adaptive Cell Decomposition (4/4)



- Close relationship with map representation (Quadtree)!


## Graph Construction: State Lattice Design (1/2)

- Enforces edge feasibility


Offline:
Motion Model


Online:
Incremental Graph Constr.

## Graph Construction: State Lattice Design (2/2)

- State lattice encodes only kinematically feasible edges




## Deterministic Graph Search

- Methods
- Breath First
- Depth First
- Dijkstra
- A* and variants
- D* and variants

| 10 | 9 | 8 | 7 | 8 |
| :---: | :---: | :---: | :---: | :---: |
| 11 | 10 |  | 4 | 7 |
|  |  |  | 5 | 6 |
| 1 | 2 |  | 4 | 5 |
| $G$ | 1 | 2 | 3 | 4 |

obstacle cell
cell with
distance value

## ADMIN

## HW4 is out

- Start early ...
- First 4 tasks are config only - build a Gazebo simulation and use ROS2 Navigation 2 to perform simple autonomy
- Task 5: Implement frontier exploration yourself $)$


## Project

- May 26 ${ }^{\text {th }}, 22: 00: \quad$ due date for the intermediate project report
- June 25 ${ }^{\text {th }}, 22: 00$ :
- June $28^{\text {th }}, 22: 00$ :
- June $28^{\text {th }}, 22: 00$ : due date for the final demo. (You are welcome to do it earlier!) due date for the final report. (You are welcome to finish it earlier!) due date for the webpage. (You are welcome to finish it earlier!)
- Schedule Project meeting with me for tomorrow, Monday or Tuesday (April 26, April 29 or April 30)!
- You contact me!
- Counts towards your "attend weekly project meeting" score!


## Project Safety!

- You work with real robotic hardware => can be dangerous!
- Don't stick your fingers where they don't belong!
- Be aware of the moving robot!
- Be careful when handling/ charging the batteries! If you are unsure, ask one of my students for help!
- Be careful with electric components.
- Be nice to the robot!


## DIJKSTRA'S ALGORITHM

## EDSGER WYBE DIJKSTRA



1930-2002
"Computer Science is no more about computers than astronomy is about telescopes."
http://www.cs.utexas.edu/~EWD/

## SINGLE-SOURCE SHORTEST PATH PROBLEM

Single-Source Shortest Path Problem - The problem of finding shortest paths from a source vertex $v$ to all other vertices in the graph.

## - Graph

- Set of vertices and edges
- Vertex:
- Place in the graph; connected by:
- Edge: connecting two vertices
- Directed or undirected (undirected in Dijkstra's Algorithm)
- Edges can have weight/ distance assigned



## Dijkstra's Algorithm

- Assign all vertices infinite distance to goal
- Assign 0 to distance from start
- Add all vertices to the queue
- While the queue is not empty:
- Select vertex with smallest distance and remove it from the queue
- Visit all neighbor vertices of that vertex,
- calculate their distance and
- update their (the neighbors) distance if the new distance is smaller


## Dijkstra's Algorithm - Pseudocode

```
dist[s]}\leftarrow 
for all v}\inV-{s
    do dist[v] }\leftarrow
S}\leftarrow
Q}\leftarrow
while Q =\varnothing
do }\textrm{u}\leftarrow\mathrm{ mindistance(Q, dist)
    S\leftarrowSU{u}
    for all v }\in\mathrm{ neighbors[u]
        do if dist[v] > dist[u] + w(u,v)
                then d[v]\leftarrowd[u]+w(u,v)
        (if desired, add traceback code)
return dist
(distance to source vertex is zero)
(set all other distances to infinity)
(S, the set of visited vertices is initially empty)
(Q, the queue initially contains all vertices)
(while the queue is not empty)
(select the element of Q with the min. distance)
(add u to list of visited vertices)
```

$\operatorname{dist}[\mathrm{s}] \leftarrow \mathrm{o}$
for all $\mathrm{v} \in \mathrm{V}-\{\mathrm{s}\}$ do dist[v] $\leftarrow \infty$
while $\mathrm{Q} \neq \varnothing$
do $\mathrm{u} \leftarrow$ mindistance( Q , dist) $\mathrm{S} \leftarrow \mathrm{S} \cup\{\mathrm{u}\}$
for all $\mathrm{v} \in$ neighbors[u]

```
(if desired, add traceback code)
return dist
```

Dijkstra Example

## Initialize:

$$
Q: \begin{array}{llllll}
A & B & C & D & E \\
\hline 0 & \infty & \infty & \infty & \infty
\end{array}
$$



$$
S:\{ \}
$$

Dijkstra Example

$$
Q: \begin{array}{lllll}
A & B & C & D & E \\
\hline 0 & \infty & \infty & \infty & \infty
\end{array}
$$



Dijkstra Example


Dijkstra Example


Dijkstra Example


Dijkstra Example


Dijkstra Example


Dijkstra Example


Dijkstra Example


Dijkstra Example


## APPLICATIONS OF DIJKSTRA'S ALGORITHM

## - Navigation Systems <br> - Internet Routing



Router A


## Dijkstra's Algorithm for Path Planning: Topological Maps

- Topological Map:
- Places (vertices) in the environment (red dots)
- Paths (edges) between them (blue lines)
- Length of path = weight of edge
- => Apply Dijkstra's Algorithm to find path from start vertex to goal vertex



## Dijkstra's Algorithm for Path Planning: Grid Maps

- Graph:
- Neighboring free cells are connected:
- 4-neighborhood: up/ down/ left right
- 8-neighborhood: also diagonals
- All edges have weight 1
- Stop once goal vertex is reached
- Per vertex: save edge over which the shortest distance from start was reached => Path


## Graph Search Strategies: Breath-First Search

- Corresponds to a wavefront expansion on a 2D grid
- Breath-First: Dijkstra‘s search where all edges have weight 1

| 10 | 9 | 8 | 7 | $8^{S}$ |
| :---: | :---: | :---: | :---: | :---: |
| 11 | 10 |  |  | 7 |
|  |  |  | 5 | 6 |
| 1 | 2 |  | 4 | 5 |
| G 2 | 1 | 2 | 3 | 4 |

## Graph Search Strategies: Depth-First Search



## Graph Search Strategies: A* Search

- Similar to Dijkstra's algorithm, except that it uses a heuristic function $h(n)$
- $f(n)=g(n)+h(n)$


A*

- Developed 1986 as part of the Shakey project!
- Complexity:

Worst-case performance $\quad O(|E|)=O\left(b^{d}\right)$
Worst-case space $\quad O(|V|)=O\left(b^{d}\right)$
complexity
b: branching factor
d: depth

- Good heuristic => small branching factor


## Optimal Planning

- Dijkstra finds the optimal path
- What about $\mathrm{A}^{*}$ ?
- Find admissible heuristic h(n)
- Admissible: do not overestimate the true cost-to-go
- $A^{*}$ is optimal (finds optimal/ shortest path) if $h(n)$ is admissible for all $n$
- Admissible example: use geometric distance for $h(n)$ :
$\left.h(n)=\sqrt{\left(x_{\text {goal }}-x\right)^{2}+\left(y_{\text {goal }}-y\right)^{2}}\right)$.
- Example: heuristic $5 x$ geometric distance
- $\mathrm{h}(\mathrm{n})=0$ => Dijkstra's Algorithm
- Hierarchical planning possible e.g.: Go to the library:
- First plan how to get from SIST building to library
- Then plan how to get from entrance of library to goal room (campus level vs. library level)
- Many variants of $\mathrm{A}^{*}$ algorithms exist - with different properties
- A* as graph search: applications outside of robotics/ path planning
- Video games
- Parsing with stochastic grammars in natural language processing
- Graph on which planning is done matters!
- E.g.: Grid map; Pose Graph; Topological Graph; Open Street Map; Lattice Graphs; ...


3D Path planned with A* in Area Graph of 2 stories of SIST \& SEM buildings

## Graph Search Strategies: D* Search

- Similar to A* search, except that the search starts from the goal outward
- $\mathrm{f}(\mathrm{n})=\mathrm{g}(\mathrm{n})+\varepsilon \mathrm{h}(\mathrm{n})$
- First pass is identical to $A^{*}$
- Subsequent passes reuse information from previous searches

$\epsilon=1.0$

$\epsilon=1.0$

$\epsilon=1.0$


## Graph Search Strategies: Randomized Search

- Most popular version is the rapidly exploring random tree (RRT)
- Well suited for high-dimensional search spaces
- Often produces highly suboptimal solutions


45 iterations


2345 iterations

## Why are RRT's rapidly exploring?

The probability of a node to be selected for expansion is proportional to the area of its Voronoi region


目

$+2+2+2+20+2$

## ritstar planner.cpp

Weat lizhi@lizhi-HP-EliteBook-8460w: ~/download codes/rrtstar_planner goal: 5.969637 .10589


Findin Found. Total paths 1 Inding Optimal Path
INFO] [1497922923.557025739, 592.300000000]: Got new plan INFO] [1497922923.957553192, 592.700000000]: Goal reached
C[rviz-13] killing on exit
[amcl-12] killing on exit
[map_server-11] killing on exit
[move_base-10] killing on exit
[kobuki_safety_controller-9] killing on exit
[navigation_velocity_smoother-8] killing on exit
onctritstarplan $h$
[cmd_vel_mux-7] killing on exit
[mobile_base_nodelet_manager-6] killing on exit
[joint_state_publisher-5] killing on exit
[diagnostic_aggregator-4] killing on exit
[robot state_publisher-3] killing on exit
[robot_state_publisher-3] killing on exit
[stageros-2] killing on exit
[rosout-1] killing on exit
[master] killing on exit
[master] killing on exit
utting down processing monitor.
shutting down processing monitor complete
done
lizhi@lizhi-HP-EliteBook-8460w:~/download codes/rrtstar_planner\$
I

## 




## teb_local_planner

An optimal trajectory planner for mobile robots based on Timed-Elastic-Bands

## General Control Scheme for Mobile Robot Systems



- Autonomous mobile robots move around in the environment. Therefore ALL of them:
- They need to know where they are.
- They need to know where their goal is.
- They need to know how to get there.
- Different levels:
- Control:
- How much power to the motors to move in that direction, reach desired speed
- Navigation:
- Avoid obstacles
- Classify the terrain in front of you
- Predict the behavior (motion) of other agents (humans, robots, animals, machines)
- Planning:
- Long distance path planning
- What is the way, optimize for certain parameters


## Navigation, Motion \& Motor Control



## Control Hierarchy

- Assume we have a goal pose



## ROS2

Navigation 2
navigation.ros.org/

More infos next class!


