

# Discussion 10

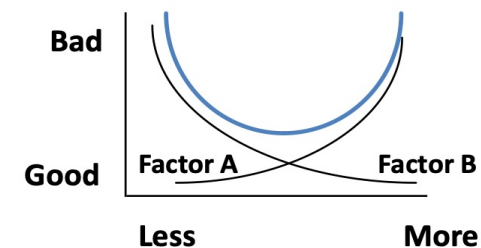
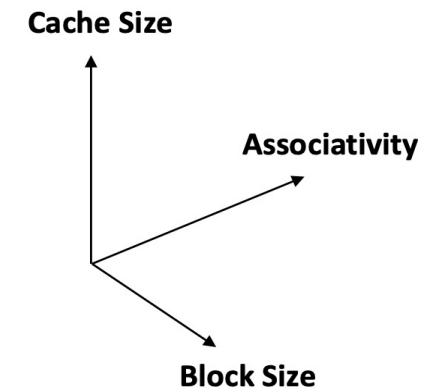
## Cache

Zheng yuting

# Review

## In Conclusion, Cache Design Space

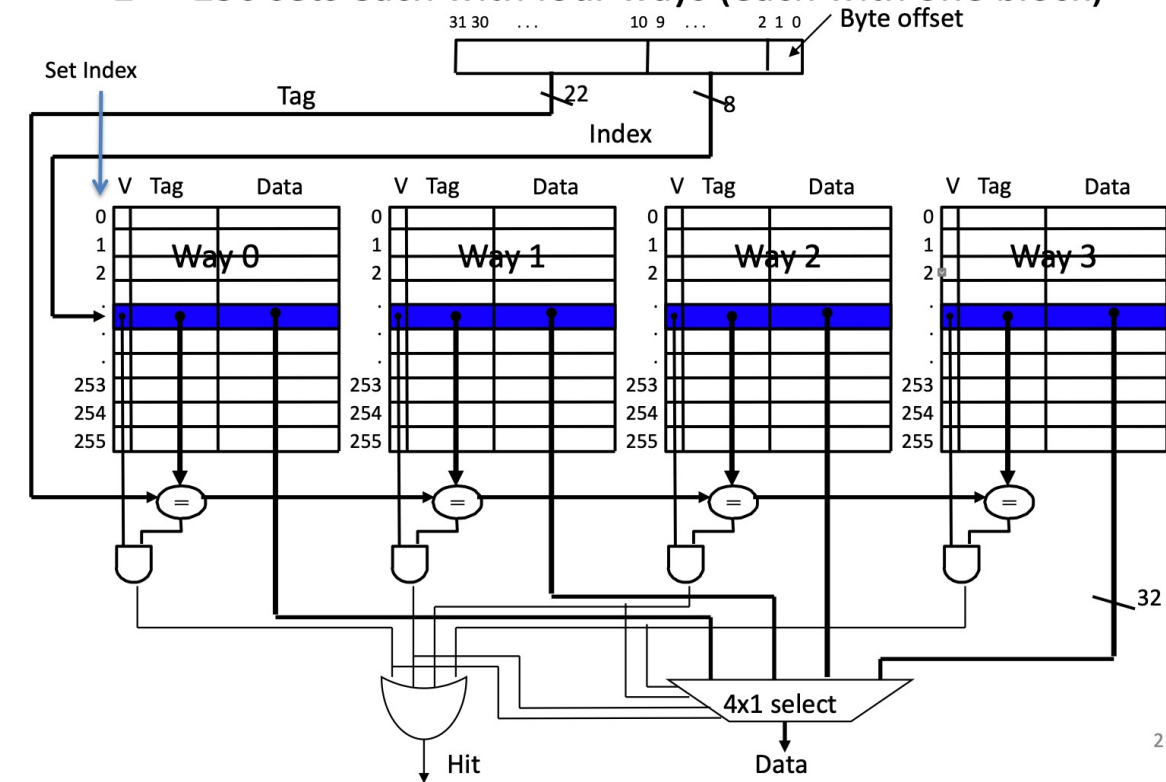
- Several interacting dimensions
  - Cache size
  - Block size
  - Associativity
  - Replacement policy
  - Write-through vs. write-back
  - Write-allocation
- Optimal choice is a compromise
  - Depends on access characteristics
    - Workload
    - Use (I-cache, D-cache)
  - Depends on technology / cost
- Simplicity often wins



# Review

## Four-Way Set-Associative Cache

- $2^8 = 256$  sets each with four ways (each with one block)



# Review

## Different Organizations of an Eight-Block Cache

**One-way set associative  
(direct mapped)**

Block	Tag	Data
0		
1		
2		
3		
4		
5		
6		
7		

**Two-way set associative**

Set	Tag	Data	Tag	Data
0				
1				
2				
3				

**Four-way set associative**

Set	Tag	Data	Tag	Data	Tag	Data	Tag	Data
0								
1								

**Eight-way set associative (fully associative)**

Tag	Data	Tag	Data	Tag	Data	Tag	Data	Tag	Data	Tag	Data	Tag	Data	Tag	Data

Total size of \$ in blocks is equal to *number of sets* × *associativity*. For fixed \$ size and fixed block size, increasing associativity decreases number of sets while increasing number of elements per set. With eight blocks, an 8-way set-associative \$ is same as a fully associative \$.

# Cache Calculation

For an  $N$ -way associative cache

$$N * \# \text{ sets} = \# \text{ blocks}$$

$$\text{cache size} = \text{block size} * \text{num blocks}$$

# Cache Problem From Previous Exam

1. An 8-way set-associative cache's total size is 4096 Bytes, the block size is 32 Bytes.  
Calculate the *index* and *tag* fields length.

# Cache Problem From Previous Exam

1. An 8-way set-associative cache's total size is 4096 Bytes, the block size is 32 Bytes.  
Calculate the *index* and *tag* fields length. 32bit machine

1. —index—:  $\log_2 \# \text{ sets} = \log_2 4096 / (32 * 8) = 4$   
—tag—:  $32 - |index| - |offset| = 32 - 4 - \log_2 32 = 23.$

# Cache Problem From Previous Exam

TAG : 22, Set Index:6, Block offset: 4, Direct mapped cache

2. Given the following C source code, what is the hit rate? Assume C processes expressions left-to-right.

```
1  #define LEN 2048
2
3  int ARRAY[LEN];
4  int main() {
5      for (int i = 0; i < LEN - 256; i+=256) {
6          ARRAY[i] = ARRAY[i] + ARRAY[i+1] + ARRAY[i+256];
7          ARRAY[i] += 10;
8      }
9  }
```

Hit rate : \_\_\_\_\_



```
#define LEN 2048

int ARRAY[LEN];
int main() {
    for (int i = 0; i < LEN - 256; i+=256) {
        ARRAY[i] = ARRAY[i] + ARRAY[i+1] + ARRAY[i+256];
        ARRAY[i] += 10;
    }
}
```

Hit rate : \_\_\_\_\_

**Solution:** 50%  
Every iteration it's  
ARRAY[i] read MISS  
ARRAY[i+1] read HIT  
ARRAY[i+256] read MISS  
ARRAY[i] write MISS (conflict! - same cache line as i+256!)  
ARRAY[i] read HIT  
ARRAY[i] write HIT  
3 MISSES, 3 HITS. 50% hit rate.

# 3Cs

1. Compulsory: First time you ask the cache for a certain block. A miss that must occur when you first bring in a block. Reduce compulsory misses by having a longer cache lines (bigger blocks), which bring in the surrounding addresses along with our requested data. Can also pre-fetch blocks beforehand using a hardware prefetcher (a special circuit that tries to guess the next few blocks that you will want).
2. Conflict: Occurs if, hypothetically, you went through the ENTIRE string of accesses with a fully associative cache and wouldn't have missed for that specific access. Increasing the associativity or improving the replacement policy would remove the miss.
3. Capacity: The only way to remove the miss is to increase the cache capacity, as even with a fully associative cache, we had to kick a block out at some point.

# Cache Problem From Previous Exam

(c) This section involves several single choice questions. **Choose the best-fit choice** for every underlined space.

1. For a cache with fixed size and associativity, enlarging the block size will \_\_\_\_\_ miss and \_\_\_\_\_ miss.

- |                         |                         |
|-------------------------|-------------------------|
| (a) increase compulsory | (b) decrease compulsory |
| (c) increase capacity   | (d) decrease capacity   |
| (e) increase conflict   | (f) decrease conflict   |

2. For a cache with fixed block size and number of sets, increasing associativity will \_\_\_\_\_ miss and \_\_\_\_\_ miss.

- |                         |                         |
|-------------------------|-------------------------|
| (a) increase compulsory | (b) decrease compulsory |
| (c) increase capacity   | (d) decrease capacity   |
| (e) increase conflict   | (f) decrease conflict   |

3. For a cache with fixed block size and length of the *tag* field, increasing associativity will \_\_\_\_\_ and \_\_\_\_\_.

- |                           |                           |
|---------------------------|---------------------------|
| (a) increase hit time     | (b) decrease hit time     |
| (c) increase miss rate    | (d) decrease miss rate    |
| (e) increase miss penalty | (f) decrease miss penalty |

1. For a cache with fixed size and associativity, enlarging the block size will \_\_\_\_ miss and \_\_\_\_ miss.  

(a) increase compulsory	(b) decrease compulsory
(c) increase capacity	(d) decrease capacity
(e) increase conflict	(f) decrease conflict
  
2. For a cache with fixed block size and number of sets, increasing associativity will \_\_\_\_ miss and \_\_\_\_ miss.  

(a) increase compulsory	(b) decrease compulsory
(c) increase capacity	(d) decrease capacity
(e) increase conflict	(f) decrease conflict
  
3. For a cache with fixed block size and length of the *tag* field, increasing associativity will \_\_\_\_ and \_\_\_\_.  

(a) increase hit time	(b) decrease hit time
(c) increase miss rate	(d) decrease miss rate
(e) increase miss penalty	(f) decrease miss penalty

**Solution:**

1. (b), (e)      (The order doesn't matter)
2. (d), (f)      (The order doesn't matter)
3. (a), (d)      (The order doesn't matter)

# Midterm Preparation

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# Q&A

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THANKS FOR YOUR ATTENDANCE